



IRP 4.3: Advanced Monitoring and Controls of the Electrical Distribution Network

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This project has received funding from the European Union's Horizon 2020 research and innovation programme under Marie Skłodowska-Curie grant agreement No 675318

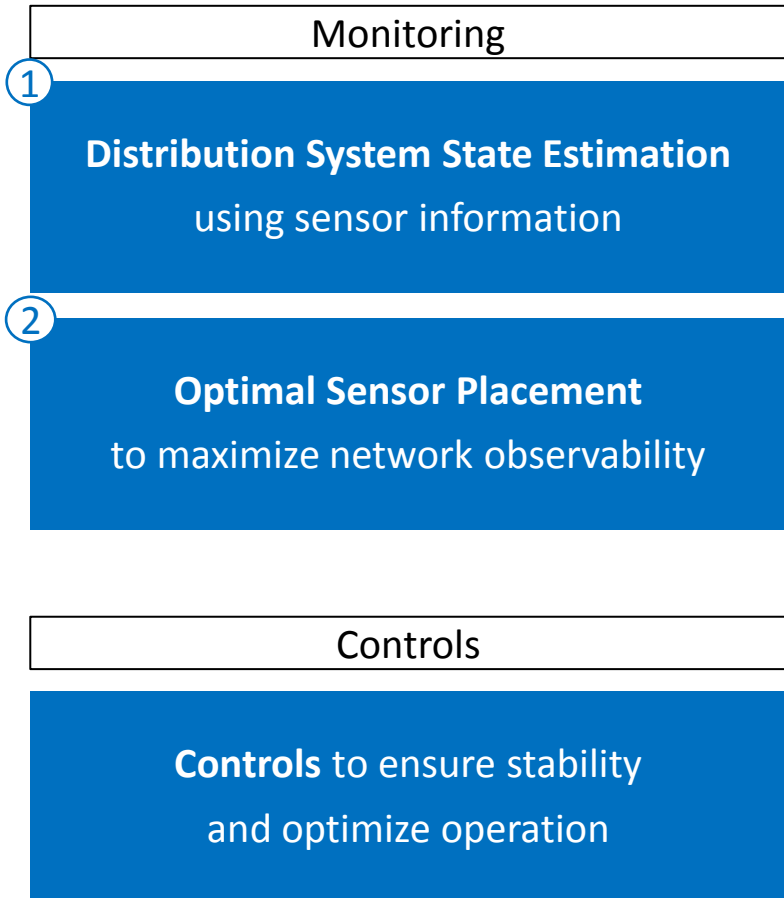


My research topic

Problems

- Large number of buses, but **few measurements available + expensive** to install all sensors. Not observable
 - **Coupled 3-phases** and **unbalanced loads**. More complex power flow
 - Simple load allocation methods based on **estimations produce large inaccuracies**
- **Not prepared** for Distributed Energy Generation (PVs, EVs, Batteries, etc.)

PhD



① Two Step State Estimation

Problem statement:

- **State Estimation** = network voltages V estimation using measurements $z = h(V) + noise$:
$$V = \operatorname{argmin}_V (z - h(V))^T \Sigma_{noise}^{-1} (z - h(V))$$
- **Mixed measurements** z :
 1. Load estimations (load forecast, installed/contracted capacity)
 2. Real-time measurements (Smart Meter, Phasor Measurement Units)

Two Step State Estimation: split problem (extending Schenato et al., 2014¹):

Prior Solution	<ul style="list-style-type: none">▪ Solve the Power Flow problem using load pseudo-measurements / estimations: $V_{prior} = PowerFlow(S_{pseudo})$	Offline Iterative High computational cost
Posterior Update	<ul style="list-style-type: none">▪ Improve prior solution using real-time measurements $z_{RT} = C_{RT}(V) + noise$: $V_{post} = V_{prior} + K(z_{RT} - C_{RT}(V_{prior}))$▪ Optimal Bayesian gain K w.r.t. trace of estimation covariance: $K = \operatorname{argmin} \operatorname{tr}(\Sigma_{post})$	Real-time / Online No iterations Low computational cost

¹ L. Schenato, G. Barchi, D. Macii, R. Arghandeh, K. Poolla, and A. Von Meier, "Bayesian linear state estimation using smart meters and pmu measurements in distribution grids," in 2014 IEEE International Conference on Smart Grid Communications (SmartGridComm).

② Optimal sensors: definition

Problem statement:

- **Sensor optimal allocation** = choose optimal number and location of sensors with budget B w.r.t. metric f :

$$x = \operatorname{argmin}_x \left(f(\Sigma_{post}(x)) \right) \text{ s. t. } \sum_i c_i x_i \leq B, x_i \in \{0,1\}$$

where $x_i = 1$ (0) if sensor i is (not) selected and $c_i \geq 0 \forall i$ is the cost of installing it

- Combinatorial optimization \rightarrow NP-hard problem

Metrics	$f =$	Covariance eigenvalues	Convex	Gradient expression	Linear	Super-modular ¹
A-opt.	$\operatorname{tr}(\Sigma_{post})$	Sum	✓	✓		
D-opt.	$\log(\det(\Sigma_{post}))$	Logarithm of product	✓	✓		✓
E-opt.	$\lambda_{max}(\Sigma_{post})$	Maximum	✓			
T-opt.	$-\operatorname{tr}(\Sigma_{post}^{-1})$? (Fisher Information)	✓	✓	✓	✓

¹ Supermodular set function $f()$: For sets $X \subseteq Y \subseteq \Omega \setminus \{a\}$, then $f(Y \cup \{a\}) - f(Y) \geq f(X \cup \{a\}) - f(X)$

② Optimal sensors: approximations

Lower bounds

Convex relaxation:

- $x_i \in \{0,1\} \rightarrow x_i \in [0,1]$, convex set
- Global minimum x_{convex} for convex relaxed problem, but not feasible w.r.t. real problem

Supermodularity minimization:

1. $\tilde{f}_{greedy} = (1 - \prod_i (1 - \frac{c_i}{B}) x_{greedy,i}) f(x_{greedy})$
 2. Select best single sensors till filling budget¹:
 $\operatorname{argmin}_{i \in \mathcal{B}, i \notin X} f(\{i\})$
- Online bound x_{online}

Upper bounds

Feasible solution:

- Select best sensors of x_{convex} till filling budget¹:
 $\operatorname{argmin}_{i \in \mathcal{B}, i \notin X} x_{convex,i}$
- Feasible, yet suboptimal solution $x_{feasible}$

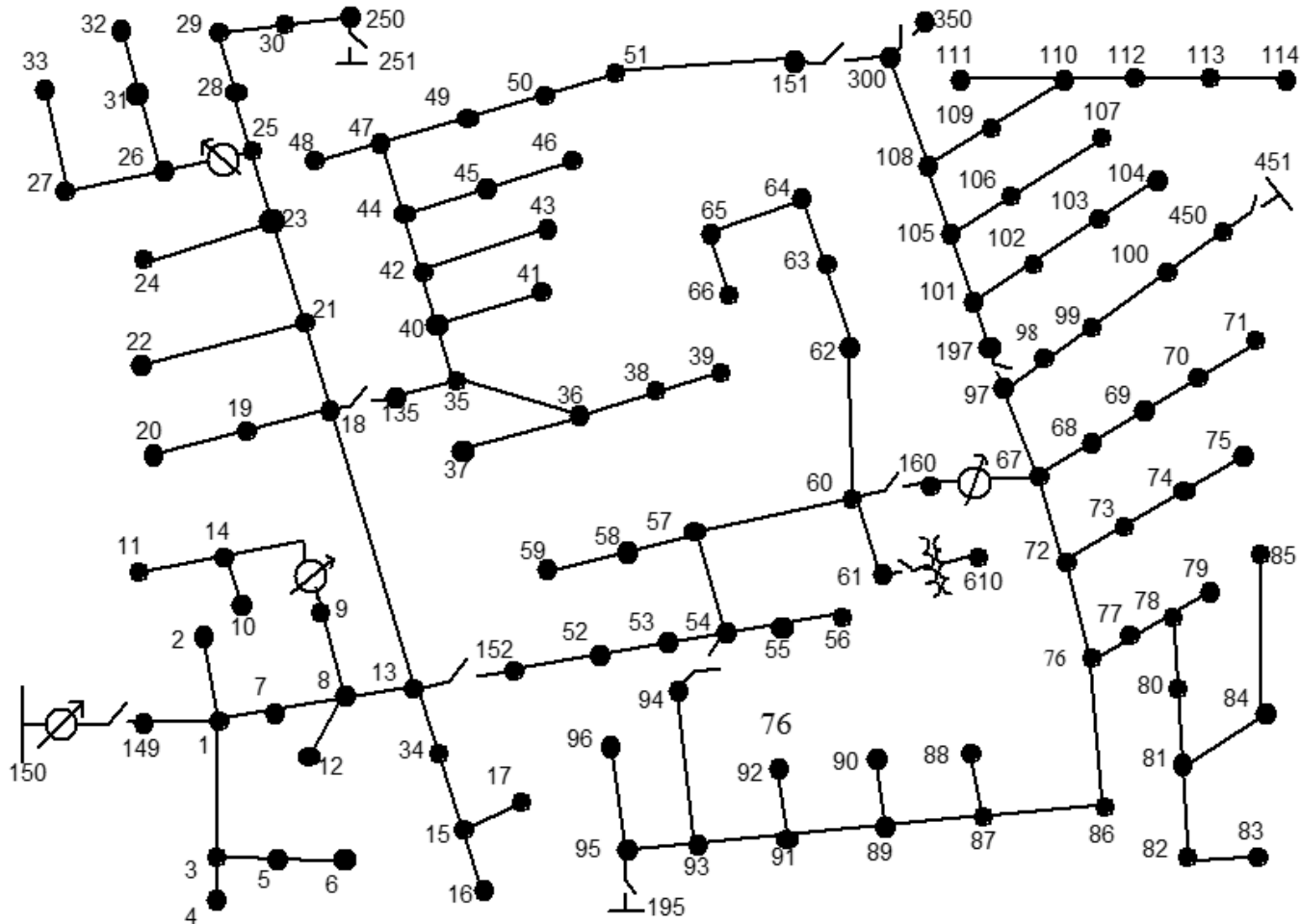
Greedy forward selection:

- Select best incremental sensor till filling budget¹:
 $\operatorname{argmin}_{i \in \mathcal{B}, i \notin X} \frac{f(X \cup \{i\}) - f(X)}{c_i}$
- Greedy feasible suboptimal solution x_{greedy}

$$\max(f(x_{convex}), f(x_{online}), \tilde{f}_{greedy}) \leq f(x_{minimum}) \leq \min(f(x_{feasible}), f(x_{greedy}))$$

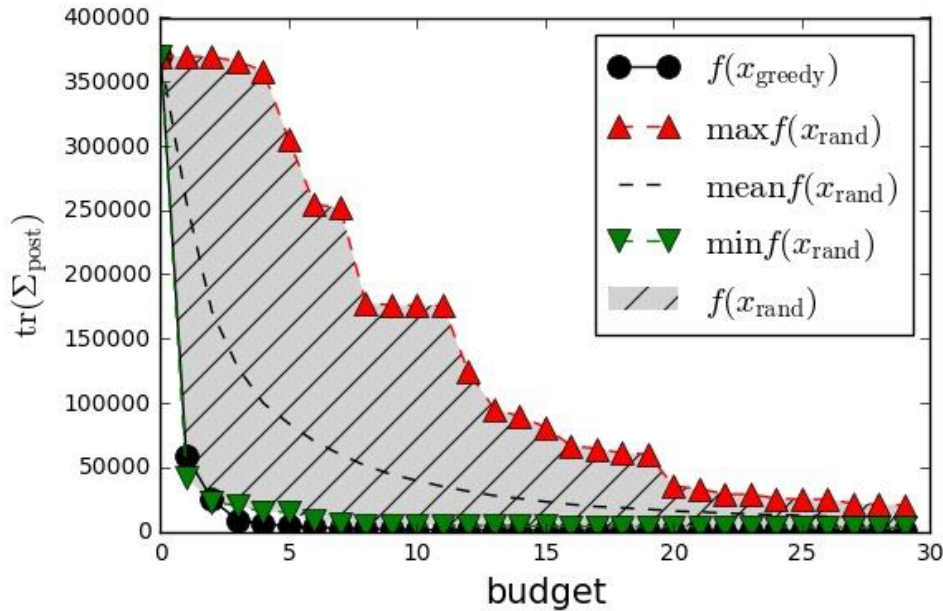
¹ Where X denotes the selected sensors by each method and $\mathcal{B} = \{i | c_i \leq \sum_{j \in X} c_j\}$ is the set of possible sensors satisfying the budget

Test Feeder: IEEE 123 Node

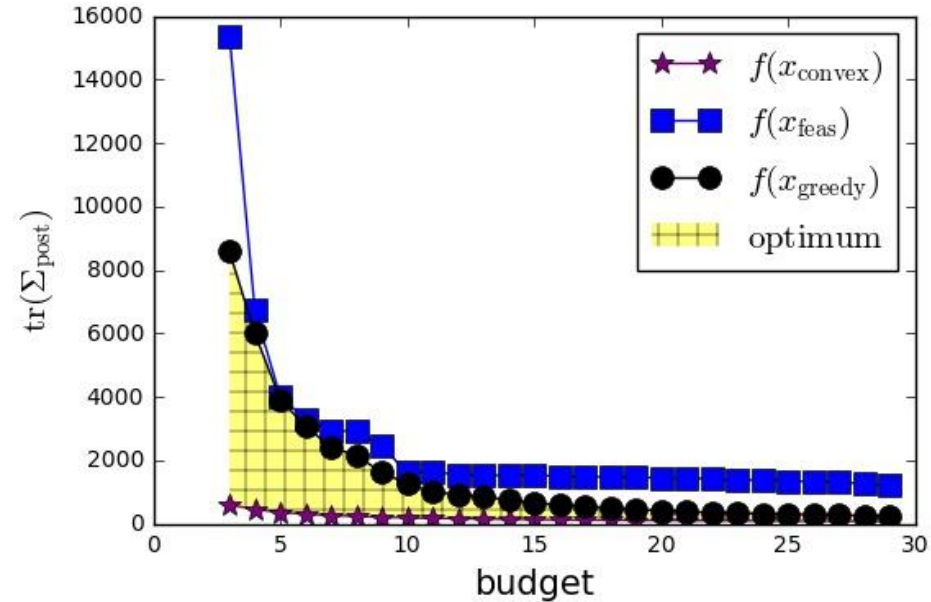


Results for A-opt.

Greedy vs Random Placement (100 samples)

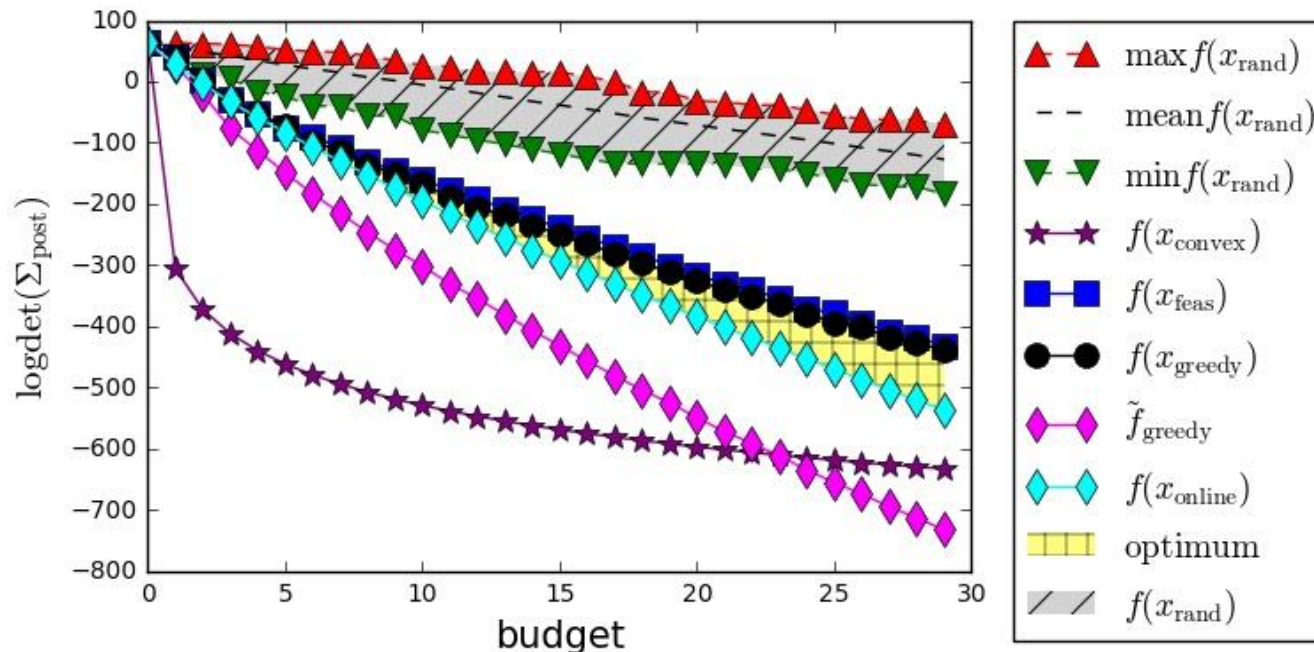


Bounds & possible optimum



- Randomized selection might be far from optimum
 - Convex relaxation bound is too optimistic for small number of sensors, but useful for many sensors
- Greedy solution is sufficiently good for large number of sensors

Results for D-opt.



- Randomized selection is far from optimum
 - Convex relaxation bound is too optimistic for small number of sensors, but useful for many sensors
 - Tight bound area for actual minimum combining all bounds
- Greedy solution is sufficiently good